

ORA: Organization Risk Analyzer*

CASOS Technical Report

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Abstract

ORA is a network analysis tool that detects risks or vulnerabilities of an organization's design structure. The design structure of an organization is the relationship among its personnel, knowledge, resource, and task entities. These entities and relationships are represented by the Meta-Matrix. Measures that take as input a Meta-Matrix are used to analyze the structural properties of an organization for potential risk. ORA contains over 50 measures which are categorized by which type of risk they detect. Measures are also organized by input requirements and by output. ORA generates formatted reports viewable on screen or in log files, and reads and writes networks in multiple data formats to be interoperable with existing network analysis packages. In addition, it has tools for graphically visualizing Meta-Matrix data and for optimizing a network's design structure. ORA uses a Java interface for ease of use, and a C++ computational backend. The current version ORA 1.2 software is available on the CASOS website <http://www.casos.cs.cmu.edu/projects/ora/software.html>.

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1. ORA Motivation and Description

ORA is a network analysis tool that detects risks or vulnerabilities of an organization's design structure. The design structure of an organization is the relationship among its personnel, knowledge, resources, and tasks entities. These entities and relationships are represented by a collection of networks called the Meta-Matrix. ORA analyzes the Meta-Matrix using measures, and reads and writes network data in multiple formats to make it interoperable with existing network analysis software.

The modeling of organizations as networks and the development of measures to examine their design structure is well developed. Even a cursory analysis of the literature reveals a wide variety of measures for assessing organizational risk and vulnerability [1] [2] [4] [6] [9] [10]. Such measures vary dramatically in the detail and type of data needed to determine that measure. They span from the assessment of critical employees, to the tendency to group think, to the potential for adaptability. In fact, it is possible to provide a suite of measures and metrics that capture both the organizational design and the possible changes in that design that are likely to result in group think, error cascades, and IP loss [7].

Given the high potential number of vulnerabilities and risks, what is needed is a framework for evaluating this set of metrics, assessing the value of existing metrics, locating gaps in the existing metrics, developing new metrics as needed, and so providing a more comprehensive guide to which metrics to use when. ORA has been designed to provide this framework.

A large number of metrics for assessing organizational vulnerability and design have been assessed and over 50 of them are now incorporated in ORA. As metrics are incorporated, if they cannot handle binary data, then we are developing a non-binary form. Help is provided for each measure that describes the measure definition and formula, input data constraints, and computational complexity. Fastest known algorithms are incorporated, employing sparse and non-sparse matrix techniques. Further, the number of nodes of any one type – personnel, knowledge, resources, tasks etc. in ORA is limited only by machine memory and processor speed. All measures are based on the Meta-Matrix and take into account the relations among personnel, knowledge, resources and tasks. These measures are based on work in social networks, operations research, organization theory, knowledge management, and task management. Where possible, metrics are normalized to be within 0 and 1 to provide a consistent framework.

ORA can be used to do a risk audit for the organization of its individual and organization risks. Such risks include, but are not limited to, tendency to groupthink, overlook of information, communication barriers, and critical employees. It evaluates potential organizational risks based upon underlying social, knowledge, resource, and task networks. This tool takes the Meta-Matrix data at a particular point in time and calculates a series of metrics assessing the team's design, particularly the command and control structure, and the associated organizational risks. ORA has been used to assess risk in various organizational and government settings including NASA, nursing hospitals, and joint task force settings.

2. Input

2.1 The Organization as Meta-Matrix

The main unit of input in ORA is the organization. An organization can be modeled and characterized as a set of interlocked networks connecting entities such as people, knowledge resources, tasks and groups. These interlocked networks can be represented using the Meta-Matrix conceptual framework (see [5] [7] [8]) presented in Table 1.

Table 1: Meta-Matrix Showing Networks of Relations Connecting Node Entities

| | People | Knowledge | Resources | Tasks/Projects |
|-----------------|--|---|--|--|
| People | Social Network <i>Who talks to, works with, and reports to whom</i> | Knowledge Network <i>Who knows what, has what expertise or skills</i> | Resource Network <i>Who has access to or can use which resource</i> | Assignment Network <i>Who is assigned to which task or project, who does what</i> |
| Knowledge | | Information Network <i>Connections among types of knowledge, mental models</i> | Resource Usage Requirements <i>What type of knowledge is needed to use that resource</i> | Knowledge Requirements <i>What type of knowledge is needed for that task or project</i> |
| Resources | | | Inter-operability and Co-usage Requirements <i>Connections among resources, substitutions</i> | Resource Requirements <i>What type of resources are needed for that task or project</i> |
| Tasks/ Projects | | | | Precedence and Dependencies <i>Which tasks are related to which</i> |

This Meta-Matrix serves as an integrating feature of a managerial toolkit. The Meta-Matrix serves several purposes; 1) it provides a way of conceptualizing the set of entities and relations among them that the research and associated tools will focus on; 2) it brings to the forefront the recognition that the data that is collecting will be not just the attributes of the entities (people, knowledge, resources, tasks and/or projects, and groups or teams) but also the set of relations or ties among them; 3) it provides an identification of the class of entities and relations that will be used in doing organizational design, analysis and risk evaluation; and 4) it provides a common ontology for talking about and representing organizational information.

2.1 Meta-Matrix Data Formats

To make ORA interoperable with existing network analysis software, ORA reads and writes Meta-Matrix network data in multiple formats. The networks that constitute a Meta-Matrix can be stored in separate files, with one network per file, or they can be collected into a single file.

ORA supports the DL, Extended-DL (EDL), CSV, and RAW formats for reading and writing a file containing a single type of network (for example, the type Agent x Agent). The RAW and DL formats are defined by the network analysis package UCINET [3]. DyNetML and EDL are the supported formats for representing in a single file the multiple network types of the Meta-Matrix.

2.1.1 DyNetML

DyNetML is an XML specification that represents the node entities Agent, Knowledge, Resource, and Task and the networks defined on them. DyNetML supports multiple Meta-Matrices to be within the same file, and each Meta-Matrix can have different Agent, Knowledge, Resource, and Task node sets. Because DyNetML is XML it is humanly readable. DyNetML is described more fully in the DyNetML Technical Report.

2.1.1 Extended-DL

Extended-DL (EDL) is the DL format of UCINET with two extensions. The first extension adds more header information to the DL file using two additional tokens. The two tokens identify the type of row nodes and the type of column nodes. The two tokens are ROW TYPE and COLUMN TYPE and are followed by one of the following: AGENT, KNOWLEDGE, RESOURCE, or TASK. COLUMN TYPE can be shortened to COL TYPE. Note that tokens in DL and EDL are case insensitive. The tokens allow the user to specify the network type of the data. For example, the following EDL file specifies a Knowledge Network (Agent x Knowledge):

```
ROW TYPE = AGENT
COLUMN TYPE = KNOWLEDGE
DL
NR=3 , NC=6
FORMAT = FULLMATRIX
DATA:
0 1 1 0 0 0
1 0 0 0 1 0
0 0 0 1 1 1
```

Note that the ROW TYPE and COLUMN TYPE tokens appear before the DL token, and because the tokens are optional, a valid DL file is a valid EDL file. In short, an EDL file can be created by adding the two new tokens to the beginning of an existing DL file.

The second extension in EDL is the ability to have multiple network types in a single file. Each network type section must be a complete and valid EDL file. The sections are separated by two vertical bars (‘||’). Thus, the extension simply allows files to be concatenated into a single

file. For example, the following file contains two types of networks from a Meta-Matrix, a Knowledge Network (Agent x Knowledge) and a Communication Network (Agent x Agent):

```
ROW TYPE = AGENT
COL TYPE = KNOWLEDGE
DL
NR=3, NC=6
FORMAT = FULLMATRIX
MATRIX LABELS:
"Knowledge Network"
DATA:
0 1 1 0 0 0
1 0 0 0 1 0
0 0 0 1 1 1
||
ROW TYPE = AGENT
COL TYPE = AGENT
DL
N=3
FORMAT = FULLMATRIX
MATRIX LABELS:
"Communication Network"
DATA:
0 1 1
1 0 0
0 0 0
```

To summarize, when ORA creates a Meta-Matrix as output the user can choose to save each type of matrix in the Meta-Matrix in a separate file, or to save the entire Meta-Matrix in a single file. When saving individual network types to a file, the following formats are available: EDL, DL, RAW, and CSV. When saving the entire Meta-Matrix to a single file, the following are available: EDL and DyNetML. These output formats make ORA interoperable with other network analysis packages.

3. Meta-Matrix Measures

An ORA measure is a function that takes a Meta-Matrix as input. Each measure examines a particular aspect of the mathematical structure of the Meta-Matrix. The metrics in ORA include social network, task management, and dynamic network metrics. ORA contains over 50 measures, and provides three classifications of them based on risk and vulnerability, input requirements, and type of output produced. The three classifications enable the user to quickly find a measure based on its properties.

3.1 Measure Risk Categories

The first classification divides the measures into seven categories of risk and vulnerability: Communication Risk, Critical Employee Risk, Resource Allocation Risk, Redundancy Risk, Personnel Interaction Risk, Task Risk, and Performance Risk. The measures in each category analyze the Meta-Matrix structure to detect the type of risk. A single measure can be classified into more than one risk category. Each of the seven categories is briefly discussed below.

3.1.1 Critical Employee Risk

Critical Employee Risk is the risk based on employees having exclusive knowledge, resources, or task assignments. Measures in this category assess in part: would the removal of one employee from the organization greatly affect the ability to complete tasks? Do employees tend to have exclusive access to knowledge or resources?

3.1.2 Resource Allocation Risk

Resource Allocation Risk is the risk based on how the organization's resource allocation affects its ability to complete tasks. Measures in this category assess: is agent workload evenly distributed? Do agents have access to the resources they need to complete tasks? Do agents have access to resources they do not use?

3.1.3 Communication Risk

Communication Risk is based on the level of communication and the authority structure of the organization. This category seeks to answer the following questions: are agents able to communicate when necessary to complete tasks? Is communication too centralized or decentralized? Do agents have recourse to managers to settle disputes?

3.1.4 Redundancy Risk

Redundancy Risk is the risk based on redundancy in task assignments, resource access, and knowledge access. An organization with little redundancy is more adversely affected by an agent or resource no longer being available. On the other hand, too much redundancy makes an organization inefficient.

3.1.5 Task Risk

Task Risk is the risk based on task precedence and task assignment. Measures in this category are able to evaluate the following questions: do agents have the resources to complete their tasks? are tasks highly interdependent so that the inability to perform one task prevents many other tasks from being completed?

3.1.6 Personnel Interaction Risk

Personnel Interaction Risk is the risk based on agent communication, either agents communicating who should not be, or vice-versa. Measures in this category examine the organization design structure to assess the following: are agents with similar skills interacting? Are agents with complementary skills interacting? Are there groups of agents communicating in unexpected ways? Is there a group of agents that has extensive reach in the organization, or whose removal would greatly fragment the organization.

3.1.7 Performance Risk

Performance Risk is the risk based on ability to complete tasks accurately. Measures in this category assess the following questions: is the organization able to complete all tasks? How well does the organization build consensus? How many tasks would be left undone if a single employee were selected for removal?

3.2 Measure Input Requirements

A second way that measures can be classified is according to input requirements. Measures that take as input a single matrix (a cell of the Meta-Matrix) are called Single-Cell measures; measures taking more than one cell are called Multi-Cell measures. In addition, some measures require only one matrix as input, but it need not correspond to a specific cell in the Meta-Matrix, but only a square sub-section of the Meta-Matrix; these are called Square measures. For example, all of the Centrality measures are Square measures, and as such they can take as input the AxA (Communication) matrix from the Meta-Matrix, but they could also take the entire Meta-Matrix – which is square, or the TxT matrix.

3.3 Measure Output Types

A third classification of measures is by output. A measure produces one of two types of output: graph level or node level. A graph level measure's output is associated with one or more matrices (also called graphs) from the Meta-Matrix. Graph-Level measures are always scalar valued. For example, Density is a graph level measures because it outputs a scalar value that describes a property of the input graph as a whole. The output of a Node Level measure, on the other hand, is associated with the members of a node entity. For example, the Cognitive Load measure is a Node Level measure because it produces a scalar value for each Agent node.

To summarize the three classification schemes, Table 2 classifies some of the measures available in ORA. A complete listing of measures available in ORA with descriptions, formulas, input requirements, and output data is in Appendix A: ORA Measures. Measures are listed by risk category in Appendix B: ORA Risk Measure Categories.

Table 2: Illustrative metrics categorized by Input, Risk, and Output

| Metric | Meaning | Output Level | Risk | Input Data |
|---------------------|--|--------------|-------------------------------|-------------|
| Degree Centrality | In the social network, number of others the person is connected to. | Node | Communication | Single-cell |
| Task Exclusivity | Detects agents who exclusively perform tasks. | Node | Critical Employee Performance | Single-cell |
| Cognitive Load | Measures the total amount of effort expended by each agent to do its tasks. | Node | Critical Employee | Multi-cell |
| Resource Congruence | Measures the similarity between what resources are assigned to tasks via agents, and what resources are required to do tasks. Perfect congruence occurs when agents have access to resources when and only when it is needful to complete tasks. | Graph | Resource Allocation, Task | Multi-cell |

4. Reports

ORA generates text reports from the measure analysis. A report is a predefined output data format. ORA currently produces a single report, the Risk and Vulnerability Report. This report can be saved in one of three formats: plain text, CSV, or DyNetML. The Risk and Vulnerability Report is a risk audit of an organization, which groups the measures by risk category and lists the measure values. The three formats are three different mediums for outputting the report data. The DyNetML report format contains in one file the original input Meta-Matrix together with all measures computed on the Meta-Matrix. The user selects which of the reports to generate, and ORA creates separate output files for each. Figure 1 below shows a portion of the text file format of a Risk and Vulnerability Report for a single Meta-Matrix organization.

Figure 1: Risk and Vulnerability Report File

```

=====  

==Performance Risk==  

=====  

Graph_Level_Measure      Type      Value  

-----  

  Omega, Knowledge       val       0.5000  

  Performance As Accuracy val       0.7980  

  Omega, Resource        val       0.4286  

Node_Level_Measure      Type      Value      Nodes_w/Value  

-----  

  Exclusivity, Task      min       0.0000     7 12 14  

                        max       1.0254     5  

                        avg       0.0834  

                        std       0.2445  

=====  

==Personnel Interaction Risk==  

=====  

Graph_Level_Measure      Type      Value  

-----  

  Distance weighted Reach val       1.0000  

  Fragmentation          val       1.0000  

Node_Level_Measure      Type      Value      Nodes_w/Value  

-----  

  Centrality, Betweenness min       0.0000     2 4 5 8 9 10 11 12 13 14 16  

                        max       0.4667     6  

                        avg       0.0923  

                        std       0.1608  

  Centrality, Inverse    min       0.0625     2 9 10 11  

                        max       0.1563     4 7  

                        avg       0.1265  

                        std       0.0373  

  Relative Expertise     min       0.0667     1 3 4 5 6 9 10 11 13 14 16  

                        max       0.0667     2 7 8 12 15  

                        avg       0.0667  

                        std       0.0000  

  Relative similarity    min       0.0000     4 9 14 16  

                        max       0.0583     6 10 11 13  

                        avg       0.0402  

                        std       0.0245

```

ORA can compare two Meta-Matrix organizations. The user can select any two meta-matrices and then generate a Risk and Vulnerability Report that compares the two organizations. Figure 2 displays a portion of one such report; it is similar to the single organization report, but it contains a side by side listing of the measure values for each organization followed by the percent by which the measure values differ.

Figure 2: Comparing Two Organizations: Risk and Vulnerability Report

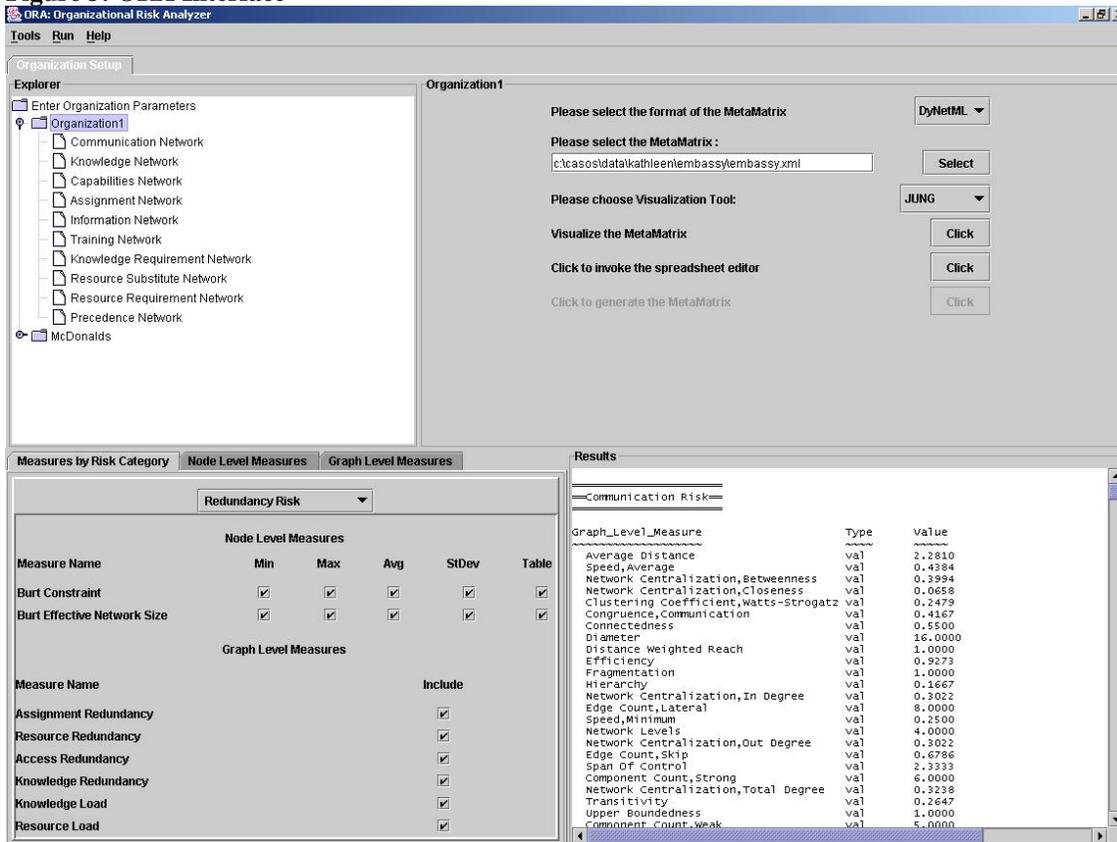
The screenshot shows a Notepad window titled 'ora-output-2orgs.txt - Notepad'. The content is a text-based report comparing two organizations across different risk categories. The report is divided into sections for 'Performance Risk' and 'Personnel Interaction Risk'. Each section contains a table with columns for 'Graph_Level_Measure', 'Type', 'organizat', 'hypoB', and '%change'. The data is presented in a structured, tabular format with dashed lines separating sections.

| ===== ==Performance Risk== ===== | | | | |
|--|------|-----------|--------|---------|
| Graph_Level_Measure | Type | organizat | hypoB | %change |
| Omega, Knowledge | val | 0.0000 | 0.0000 | |
| Performance AS Accuracy | val | 0.8292 | 0.8300 | 0 |
| Omega, Resource | val | 0.0000 | 0.0000 | |
| Node_Level_Measure | Type | organizat | hypoB | %change |
| Exclusivity, Task | min | 1.0000 | 1.0000 | 0 |
| | max | 1.0000 | 1.0000 | 0 |
| | avg | 1.0000 | 1.0000 | 0 |
| | std | 0.0000 | 0.0000 | |
| ===== ==Personnel Interaction Risk== ===== | | | | |
| Graph_Level_Measure | Type | organizat | hypoB | %change |
| Distance weighted Reach | val | 1.0000 | 1.0000 | 0 |
| Fragmentation | val | 1.0000 | 1.0000 | 0 |
| Node_Level_Measure | Type | organizat | hypoB | %change |
| Centrality, Betweenness | min | 0.0000 | 0.0000 | |
| | max | 0.4000 | 0.0000 | -100 |
| | avg | 0.1333 | 0.0000 | -100 |
| | std | 0.1491 | 0.0000 | -100 |
| Centrality, Inverse | min | 0.1667 | 0.1667 | 0 |
| | max | 0.2778 | 0.2000 | -28 |
| | avg | 0.2331 | 0.1944 | -16 |
| | std | 0.0399 | 0.0124 | -68 |
| Relative Expertise | min | 0.0000 | 0.0000 | |
| | max | 0.2000 | 0.2000 | 0 |
| | avg | 0.1667 | 0.1667 | 0 |
| | std | 0.0745 | 0.0745 | 0 |
| Relative similarity | min | 0.1000 | 0.1000 | 0 |
| | max | 0.1333 | 0.1333 | 0 |
| | avg | 0.1175 | 0.1175 | 0 |
| | std | 0.0121 | 0.0121 | 0 |

5. User Interface Components

ORA has a Java user interface for cross platform compatibility. The interface contains three main components: 1) Meta-Matrix Manager, 2) Measure Manager, and 3) Output Panel. In addition, it contains the following sub-components: the Visualizer, the Optimizer, and the Regression Tool. Each of these will be briefly described in turn. The descriptions will refer to Figure 3 below, which shows the three main components of the ORA interface.

Figure 3: ORA Interface



5.1 Meta-Matrix Manager

The unifying concept in ORA is the Meta-Matrix. The user needs to be able to manage multiple organizations, both those entered as original input and those output by the Optimizer. These are collected and managed in the Meta-Matrix Manager, which occupies the upper half of Figure 3. From this panel, the user can add organizations to the collection, rename them, and specify data files for each network of the Meta-Matrix. Most user actions in ORA require the selection of one or more organizations as input, and the available organizations are those entered into the Meta-Matrix Manager.

5.2 Measure Manager

The Measure Manager is a separate panel, seen in the lower left panel of Figure 3, that gives the user different views of the measures, according to the categories described above: Risk Category, Node Level, and Graph Level. These are all views of the same set of underlying measures, so selecting or unselecting a measure in one view is reflected in all other views.

5.3 Output Panel

The Output Panel is a text panel that gives the user immediate feedback. The Output Panel displays status information from user actions, and also displays the text content of measure reports. The Output Panel displaying a Risk Report can be seen in the lower right half of Figure 3.

5.4 Tools

The Visualizer, Optimizer, and Regression Tool are invoked from the main menu and are contained within pop-up windows. The different popup components of ORA constitute an integrated graphical user interface that has proven to be extensible and flexible.

5.4.1 Visualization

The Meta-Matrix contains multiple node entities and different types of edges. Most existing visualization packages cannot display multiple network types simultaneously, and therefore are not suitable for visualizing the Meta-Matrix. ORA contains two integrated visualization packages for displaying an entire Meta-Matrix: NetworkViz and Jung.

NetworkViz was developed at CASOS specifically for visually analyzing Meta-Matrix data, and is capable of displaying all of the networks of the Meta-Matrix simultaneously. NetworkViz can also display specific parts of the Meta-Matrix, for example, a single network, or all networks defined on one or more node entities. In this manner, the user can isolate and visualize portions of the Meta-Matrix that are of interest. If ORA has computed measures for the Meta-Matrix, then NetworkViz displays them. For example, all Node Level measures computed for a particular Agent node are displayed in a pop-up window when the node is right clicked. Figure 4 contains a sample Meta-Matrix visualization using NetworkViz.

Figure 4: NetworkViz Visualizer

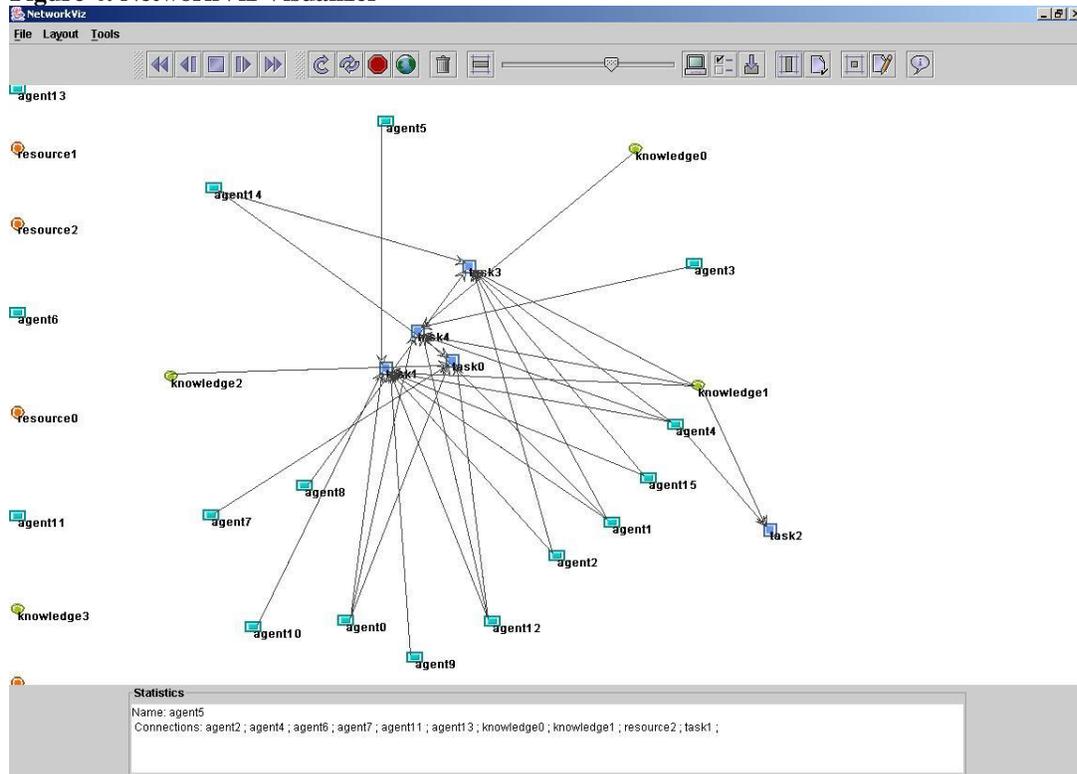
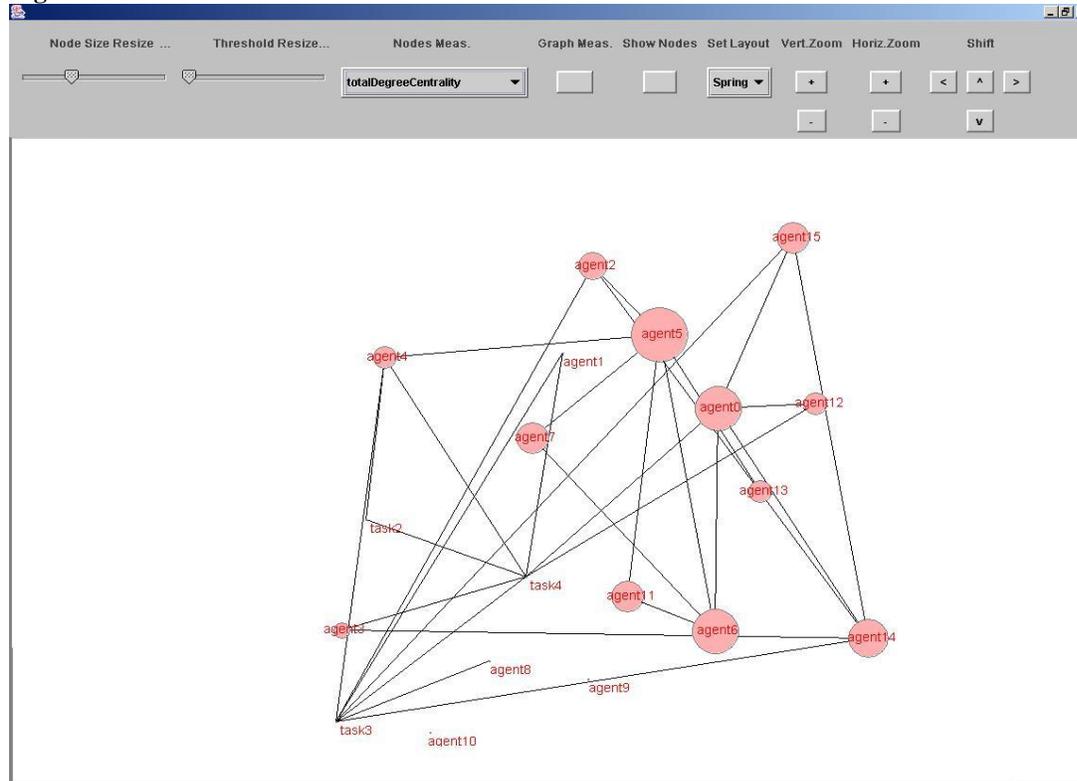


Figure 5: JUNG Visualizer



The Jung visualization package is an open-source code project from Stanford adapted and integrated into ORA. It offers three different layout algorithms, and has been customized to display Node Level and Graph Level measures computed on the Meta-Matrix. With Jung the user can choose portions of the Meta-Matrix to display based on measure values, for example, Figure 5 displays the Agent nodes with sizes proportional to their Total Degree Centrality.

5.4.2 Optimizer

Having detected the risk and vulnerabilities of an organization's design structure, the Optimizer tool in ORA allows the user to change the structure according to user specified criteria. The user selects a single measure or a linear combination of measures to be an objective function, and the Optimizer produces an organization that maximizes or minimizes the objective function by adding and removing relationships (i.e. edges) between node entities. Because the output of the Optimizer is a Meta-Matrix, it can be input to ORA for measure analysis and visualization. Details of the Optimizer can be found in the Optimizer Technical Report.

5.4.3 Regression Tool

The Regression Tool allows the user to compare two vector valued measures, plotting the two vectors in coordinate space with a linear regression line. The two vectors can be the same measure computed on two different organizations, two measures computed on the same organization, or two different measures computed on two different organizations. The plot output can be saved to a file.

6. System Requirements

ORA 1.2 is the latest version of ORA and it runs on any Windows 2000 or XP machine running on an Intel processor. The C++ back-end source code is written so as to be compatible with platforms and processors, and is being ported and tested on other platforms and processors.

7. Conclusion

ORA advances the state of the art in network analysis tools by being organized around the unifying concept of the Meta-Matrix. Measures are organized to facilitate their coherent use. In particular, they are categorized by how they measure the risk and vulnerability of an organization's design structure. ORA reads and writes in multiple data formats and is interoperable with existing network analysis software. Entire Meta-Matrices can be visualized using different layout algorithms. The integrated Optimizer adapts an organization's design structure according to user specified criteria, and the resulting organization can be input into ORA and analyzed and visualized. The computational back end employs NetStatPlus, an open source C++ library of SNA and DNA routines. The Java graphical user interface is designed for ease of use and for extensibility and flexibility as new features are added. ORA is being actively developed and tested in a wide range of contexts.

8. Future Work

Future work in ORA will address all aspects of its core functionality, including (1) managing Meta-Matrix organizations; (2) measure presentation and selection; (3) network visualization; (4) tool sub-components; and (5) generating reports.

Currently a Meta-Matrix can contain only one matrix of each network type. Thus a Meta-Matrix cannot have a Communication Network and a Friendship Network, both of type Agent x Agent. Similarly, time period data for a matrix type is not possible. Future versions will extend the Meta-Matrix Manager to allow multiple matrices of a single type.

The Measures Manager currently does not allow the user to specify the input for measures. Certain measures have predefined input, and so specifying input matrices is unnecessary. For example the Actual Knowledge Workload measure always takes as input the following matrices: Agent x Knowledge, Knowledge x Task, and Agent x Task. Other measures, for example Square measures, can operate on any square input matrix. For example, the Betweenness Centrality measure takes any square matrix as input. Currently, such measures run on a pre-determined, default matrix which is not user selectable.

Another tool currently being developed for ORA called the Matrix Tool, which displays matrix data in an editable spreadsheet window. Individual networks can be displayed, or an entire Meta-Matrix. The Matrix Tool lets the user manipulate matrices, such as performing matrix algebra, and computing the Intersection, Union, and Central Matrix of a collection of matrices. By loading a Meta-Matrix in one format and saving in another, the user can convert data from one format to another. The Matrix Tool will be included in the next release of ORA.

Finally, ORA will be extended to provide multiple report types. Currently, only the Risk and Vulnerability Report is available. Future report types will output alphabetical lists of measures, or measures categorized by node, and graph level. The Output Panel of the user interface will be extended to display multiple output files, allowing the user to quickly organize and view report files.

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Appendix A – ORA Measures

A **network** N consists of two sets of nodes, called U and V , and a set of edges $E \subset U \times V$. An element $e = (i,j)$ in E indicates a relationship or tie between nodes $i \in U$ and $j \in V$. A network where $U=V$ and therefore $E \subset V \times V$ is called **unimodal**; otherwise the network is **bimodal**. For our purposes, unimodal networks will not contain self loops, which means that $(i,i) \notin E$ for $i \in V$.

An **organization** is a collection of networks. A **measure** is a function that maps one or more networks to R^n . Measures are often scalar ($n=1$) or vector valued with $n = |V|$ or $n=|U|$.

When defining or implementing measures, a network can be represented as (1) a graph, or as (2) an adjacency matrix. To represent a *unimodal* network as a graph, let $G=(V,E)$, where V is the network's nodes, and E are the ties; *bimodal* networks will not be represented as graphs. Both unimodal and bimodal networks are represented as adjacency matrices.

Given a network $N=((U,V),E)$, define a matrix M of dimension $|U| \times |V|$, and let $M(i,j) = 1$ if $(i,j) \in E$, else let $M(i,j)=0$. Then M is called the adjacency matrix representation of network N . Unimodal networks are also called **square** networks because their adjacency matrix is square; the diagonal is zero diagonal because there are no self-loops.

Define the following sets of nodes: Agents, Knowledge, Resources, and Tasks. The following networks defined on these node sets are used throughout the documentation:

| Symbol | Node Sets | | Name |
|-----------|-----------|-----------|-------------------------------|
| | U | V | |
| A | Agent | Agent | Communication Network |
| AK | Agent | Knowledge | Knowledge Network |
| AR | Agent | Resource | Capabilities Network |
| AT | Agent | Task | Assignment Network |
| K | Knowledge | Knowledge | Information Network |
| KR | Knowledge | Resource | Training Network |
| KT | Knowledge | Task | Knowledge Requirement Network |
| R | Resource | Resource | Resource Substitute Network |
| RT | Resource | Task | Resource Requirement Network |
| T | Task | Task | Precedence Network |

The following matrix notation is used:

$|\text{Matrix}|$ = dimension of a *square* Matrix (i.e. if Matrix has dimension $r \times r$, then $|\text{Matrix}| = r$)
 $\text{Matrix}(i,j)$ = the entry in the i^{th} row and j^{th} column of Matrix
 $\text{Matrix}(i,:)$ = i^{th} row vector of Matrix
 $\text{Matrix}(:,j)$ = j^{th} column vector of Matrix
 $\text{sum}(\text{Matrix})$ = sum of the elements in Matrix (also, Matrix can be a row or column vector of Matrix)
 Matrix' = the transpose of Matrix
 $\sim\text{Matrix}$ = for binary Matrix, $\sim\text{Matrix}(i,j) = 1$ iff $\text{Matrix}(i,j) = 0$.
 $\text{Matrix}@\text{Matrix}$ = element-wise multiplication of two matrices (e.g. $C=A@B \Rightarrow C(i,j) = A(i,j)*B(i,j)$)

These mathematical terms and symbols are used:

$\text{card}(\text{Set}) = |\text{Set}|$ = the cardinality of Set
 $\text{sgn}(x) = 1$ if $x \geq 0$, and -1 otherwise
 \mathfrak{R} denotes a real number
 \mathbb{Z} denotes an integer

These graph theoretic terms are used:

$d_G(i, j)$ is the length of the shortest directed path in G from node i to node j . Note that if there is a path from i to j in G , then $1 \leq d_G(i, j) < |V|$. Therefore, let $d_G(i, j) = |V|$ if there is no path in G from i to j . Also, let $d_G(i, i) = 0$ for each $i \in V$.

The **Reachability Graph** for a square network $N=(V,E)$ is defined as follows: let $G=(V,E)$ be the graph representation for N . The Reachability Graph for N is the graph $G'=(V,E')$ where $E' = \{(i,j) \in V \times V \mid \exists \text{ directed path from } i \text{ to } j \text{ in } G\}$.

The **Underlying Network** for a network $N=(V,E)$ is defined as follows: $N'=(V,E')$ where $E' = \{(i,j) \mid (i,j) \in E \vee (j,i) \in E\}$. That is, an symmetric version of N .

| Measure | Description | Reference | Formula |
|-------------------------------|---|----------------|--|
| Access Index, Knowledge Based | <p>Boolean value which is true if an agent is the only agent who knows a piece of knowledge and who is known by exactly one other agent. The one agent known also has its KAI set to one.</p> <p>Type Node Level Input AK:binary; A:binary Output Binary</p> | Ashworth, 2003 | <p>The Knowledge Access Index (KAI) for agent i is defined as follows: let $S_i = \{s \mid AK(i, s) \wedge (\text{sum}(AK(:, s)) = 1) \wedge (\text{sum}(A(i, :)) = 1)\}$ Then $KAI_i = ((S_i \neq \emptyset) \vee (\exists j \mid S_j \neq \emptyset \wedge A(j, i) = 1))$</p> |
| Access Index, Resource Based | <p>Boolean value which is true if an agent is the only agent with access to a resource and who is known by exactly one other agent. The one agent known also has its RAI set to one.</p> <p>Type Node Level Input AR:binary; A:binary Output Binary</p> | Ashworth, 2003 | <p>The Resource Access Index (RAI) for agent i is defined identically as Knowledge Access Index, with the matrix AK replaced by AR.</p> |
| Actual Workload, Knowledge | <p>The knowledge an agent uses to perform the tasks to which it is assigned.</p> <p>Type Node Level Input AK:binary; KT:binary; AT:binary Output $\mathcal{R} \in [0,1]$</p> | Carley, 2002 | <p>Actual Workload for agent i is defined as follows:</p> $[AK * KT * AT'](i, i) / \text{sum}(KT)$ <p>Note how Potential Workload is the first matrix product.</p> |
| Actual Workload, Resource | <p>The resources an agent uses to perform the tasks to which it is assigned.</p> <p>Type Node Level Input AR:binary; RT:binary; AT:binary Output $\mathcal{R} \in [0,1]$</p> | Carley, 2002 | <p>Actual Resource Workload for agent i is identical to Actual Knowledge Workload, replacing AK with AR and KT with RT.</p> |
| Average Distance | <p>The average shortest path length between nodes, excluding infinite distances.</p> <p>Type Graph Level Input A:binary, square Output $\mathcal{R} \in [0,1]$</p> | NetStat | <p>Let $G=(V,E)$ represent a square network. Define a set S of all pairs (i,j) of nodes such that i can reach j. Then average the shortest paths.</p> <p>let $S = \{(i,j) \mid j \text{ is reachable in } G \text{ from } i\}$</p> $\sum_{(i,j) \in S} d_G(i, j)$ <p>Then, Average Distance = $\frac{\sum_{(i,j) \in S} d_G(i, j)}{ S }$.</p> |

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| <p>Centrality, Betweenness</p> | <p>The Betweenness Centrality of node v in a network is defined as: across all node pairs that have a shortest path containing v, the percentage that pass through v. This is defined for directed networks. Type Node Level Input N: square Output $\mathfrak{R} \in [0,1]$</p> | <p>Freeman, 1979</p> | <p>Let $G=(V,E)$ be the graph representation for the network. Let $n= V$, and fix a node $v \in V$. For $(u,w) \in V \times V$, let $n_G(u,w)$ be the number of geodesics in G from u to w. If $(u,w) \in E$, then set $n_G(u,w)=1$. Define the following: let $S = \{(u,w) \in V \times V \mid d_G(u,w) = d_G(u,v) + d_G(v,w)\}$ let between = $\sum_{(u,w) \in S} (n_G(u,v) * n_G(v,w)) / n_G(u,w)$ Then Betweenness Centrality of node $v = \text{between} / ((n-1)(n-2)/2)$. Note: if G is not symmetric, then between is normalized by $(n-1)(n-2)$.</p> |
| <p>Centrality, Closeness</p> | <p>The average closeness of a node to the other nodes in a network. Loosely, Closeness is the inverse of the average distance in the network between the node and all other nodes. This is defined for directed networks. Type Node Level Input N:square Output $\mathfrak{R} \in [0,1]$</p> | <p>Freeman, 1979</p> | <p>Let $G=(V,E)$ be the graph representation of the square network. Fix $v \in V$. let $\text{dist} = \sum_{i \in V} d_G(v,i)$, if every node is reachable from v Then Closeness Centrality of node $v = (V -1)/\text{dist}$. If some node is not reachable from v then the Closeness Centrality of v is V.</p> |
| <p>Centrality, Eigenvector</p> | <p>Calculates the eigenvector of the largest positive eigenvalue of the adjacency matrix representation of a square network. Type Node Level Input N:square, symmetric Output $\mathfrak{R} \in [0,1]$</p> | <p>Bonacich P, 1972</p> | <p>Calculates the eigenvector of the largest positive eigenvalue of the adjacency matrix representation of a square network. A Jacobi method is used to compute the eigenvalues and vectors.</p> |
| <p>Centrality, In Degree</p> | <p>The In Degree Centrality of a node in a unimodal network is its normalized in-degree. Type Node Level Input N:square Output $\mathfrak{R} \in [0,1]$</p> | <p>Wasserman and Faust, 1994</p> | <p>Let $G=(V,E)$ be the graph representation of a square network and fix a node v. let $\text{deg} = \text{card}\{u \in V \mid (u,v) \in E\}$, this is the in-degree of node v. The In Degree Centrality of node $v = \text{deg} / (V -1)$</p> |

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| Centrality, Information | Calculate the Stephenson and Zelen information centrality measure for each node. Type Node Level Input N:square, symmetric Output $\mathcal{R} \in [0,1]$ | Wasserman and Faust, 1994 (pg 195) | Calculates the measure described on pg 195-6 of Wasserman and Faust. Nodes with 0 degree are first removed from the network, and the measure computed on the resulting sub-graph. The removed nodes are given centrality value 0. |
| Centrality, Inverse Closeness | The average closeness of a node to the other nodes in a network. Inverse Closeness is the sum of the inverse distances between a node and all other nodes. This is defined for directed networks. Type Node Level Input N:square Output $\mathcal{R} \in [0,1]$ | Wasserman and Faust, 1994 (pg 195) | Let $G=(V,E)$ be the graph representation of the square network. Fix $v \in V$. let $\text{dist} = \sum_{i \in V} \frac{1}{d_G(v,i)}$, where $\frac{1}{d_G(i,i)} = 0$ and $\frac{1}{d_G(v,i)} = 0$ if i is not reachable from v . Then Inverse Closeness Centrality of node $v = \text{dist}/(V -1)$. |
| Centrality, Out Degree | The Out Degree Centrality of a node in a square network is its normalized out-degree. Type Node Level Input N:square Output $\mathcal{R} \in [0,1]$ | Wasserman and Faust, 1994 | Let $G=(V,E)$ be the graph representation of a square network and fix a node v . let $\text{deg} = \text{card}\{u \in V \mid (v,u) \in E\}$, this is the out-degree of node v . The Out Degree Centrality of node $v = \text{deg} / (V -1)$ |
| Centrality, Total Degree | The Total Degree Centrality of a node in a square network is its normalized in plus out degree. Type Node Level Input N:square, undirected Output $\mathcal{R} \in [0,1]$ | Wasserman and Faust, 1994 (pg 199) | Let $G=(V,E)$ be the graph representation of a square network and fix a node v . let $\text{deg} = \text{card}\{u \in V \mid (v,u) \in E \vee (u,v) \in E\}$, this is the total degree of node v . The Total Degree Centrality of node $v = \text{deg} / 2*(V -1)$ |

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| <p>Clustering Coefficient, Watts-Strogatz</p> | <p>Measures the degree of clustering in a network by averaging the clustering coefficient of each node i, defined as the ratio of the number of triangles connected to i to the number of triples centered at i. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$</p> | <p>Watts and Strogatz, 1998</p> | <p>let $G=(V,E)$ be the graph representation of a square network. For each node $v \in V$ define the following: let $in_v = \{i \in V \mid (i,v) \in E\}$ let $out_v = \{i \in V \mid (v,i) \in E\}$ let $inconnect_v = \{(i,j) \in E \mid i,j \in in_v\}$ let $outconnect_v = \{(i,j) \in E \mid i,j \in out_v\}$ Then compute for each node $v \in V$ its Clustering Coefficient cc_v using (1) in-degree, (2) out-degree, or (3) total degree. (1) let $cc_v = \frac{ inconnect_v }{ in_v ^2 - in_v }$, if $in_v > 1$, else $cc_v = 0$. (2) let $cc_v = \frac{ outconnect_v }{ out_v ^2 - out_v }$, if $out_v > 1$, else $cc_v = 0$. (3) let $cc_v = \frac{1}{2}(case(1) + case(2))$ Then Clustering Coefficient for the graph = $\left(\sum_{v \in V} cc_v \right) / V$</p> |
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| <p>Cognitive Load</p> | <p>Measures the total amount of effort expended by each agent to do its tasks.</p> <p>Note: Cognitive Load is defined if one or both of the following pairs of networks exists: {AR,RT}, {AK,KT}.</p> <p>Type Node Level Input A:binary; AT:binary; [AR:binary; RT:binary]; [AK:binary; KT:binary] Output $\mathfrak{R} \in [0,1]$</p> | <p>Carley, 2002</p> | <p>The Cognitive Load for agent i is defined as follows: let $ATR = AT*RT'$ let $ATA = AT*AT'$ let $x_1 = \#$ of agents that agent i interacts with / total # of agents</p> $= \left(\sum_{j \neq i} A(i, j) \right) / (A - 1)$ <p>let $x_2 = \#$ of tasks agent i is assigned to / total # of tasks $= \text{sum}(AT(i,:)) / T$ let $x_3 = \text{sum of } \# \text{ agents who do the same tasks as agent i} / (\text{total } \# \text{ tasks} * \text{total } \# \text{ agents})$</p> $= \left(\sum_{j \neq i} ATA(i, j) \right) / (A - 1)(T)$ <p>Note that x_4, x_5, x_6 depend upon networks AR and RT; if the networks AK and KT exist, then three analogous terms for knowledge are computed and averaged. If only AK and KT exist, then only they are used.</p> let $x_4 = \#$ of resources agent i manages / total # of resources $= \text{sum}(AR(i,:)) / R $ let $x_5 = \text{sum of } \# \text{ resources agent i needs to do all its tasks} / (\text{total } \# \text{ tasks} * \text{total } \# \text{ resources})$ $= \text{sum}(ATR(i,:)) / (T * R)$ let $x_6 = \text{sum of negotiation needs agent i must do for each task} / \text{total possible negotiations}$ $= \left(\sum_j (AR(i, j) > 0 \neq ATR(i, j) > 0) \right) / (R T)$ <p>Then Cognitive Load for agent i = $(x_1 + x_2 + x_3 + x_4 + x_5 + x_6) / 6$</p> |
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| Communication | Measures the communication need of agents to complete their assigned tasks. Type Node Level Input A:binary; AT:binary; AR:binary; RT:binary, T:binary Output $\mathcal{R} \in [0,1]$ | Carley, 2003 | Communication uses the concepts from Communication Congruence: Handoff, Co-Assignment, and Negotiation. let H, C, and N be defined as in Communication Congruence. let $M(i,j) = [A + (H+H') + C + (N+N')](i,j) > 0$, and $M(i,i) = 0$ Note that the transpose of H and N is used to make the communication reciprocal. let $d = \text{sum}(M(i,:))$ let $d = d / (A -1)$, normalizing d to be in [0,1] Then Communication for agent i is d. |
| Communicative Need | Measures the percentage of reciprocal edges in a network. Type Graph Level Input N:square Output $\mathcal{R} \in [0,1]$ | Carley, 2002 | Let $G = (V,E)$ represent a square network: Then the Communicative Need = (Reciprocal Edge Count of G) / E |
| Component Count, Strong | The number of strongly connected components in a network. Type Graph Level Input N:square Output $Z \in [0, V]$ | Wasserman and Faust, 1994 (pg 109) | Given a square network represented by a graph $G=(V,E)$, the Strong Component Count is the number of strongly connected components in G. This is computed directly on G, whether or not G is directed. |
| Component Count, Weak | The number of weakly connected components in a network. Type Graph Level Input N:square, symmetric Output $Z \in [0, V]$ | Wasserman and Faust, 1994 (pg 109) | Given a square, symmetric network represented by a graph $G=(V,E)$, the Weak Component Count is the number of connected components in G. Such components are called “weak” because the graph G is undirected. |

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| <p>Congruence, Communication</p> | <p>Measures to what extent the agents communicate when and only when it is needful to complete tasks. Perfect congruence requires reciprocal communication.</p> <p>Type Graph Level Input A:binary; AT:binary; AR:binary; RT:binary, T:binary Output $\mathcal{R} \in [0,1]$</p> | <p>Carley, 2002</p> | <p>Communication Congruence = 1 iff agents communicate when and only when it is needful to complete their tasks. There are three task related reasons when agents i and j need to communicate:</p> <ul style="list-style-type: none"> (a) <i>Handoff</i>: if i is assigned to a task s and j is assigned to a task t and s directly precedes task t (b) <i>Co-Assignment</i>: if i is assigned to a task s and j is also assigned to s (c) <i>Negotiation</i>: if i is assigned to a task s and j is not, and there is a resource r to which agents assigned to s have no access but j does. <p>The three cases are computed as follows:</p> <ul style="list-style-type: none"> (a) let $H = AT * T' * AT'$ (b) let $C = AT * AT'$ (c) let $N = AT * Z * AR'$, where $Z(t,r) = [AT' * AR - RT'](t,r) < 0$ <p>Note that C is always symmetric, but not necessarily H and N.</p> <p>let $Q(i,j) = [(H+H') + C + (N+N')](i,j) > 0$.</p> <p>Communication Congruence requires reciprocal communication, explaining the transposes of H and N to make them symmetric.</p> <p>let d = hamming distance between Q and A, which measures the degree to which communication differs from that which is needed to do tasks. The maximum value for d is $d_{max} = A * (A - 1)$</p> <p>Then Communication Congruence = $1 - (d / d_{max})$, which is in $[0,1]$.</p> |
| <p>Congruence, Knowledge</p> | <p>Measures the similarity between what knowledge is assigned to tasks via agents, and what knowledge is required to do tasks. Perfect congruence occurs when agents have knowledge when and only when it is needful to complete tasks.</p> <p>Type Graph Level Input AK:binary; AT:binary; KT:binary Output $\mathcal{R} \in [0,1]$</p> | <p>Carley, 2002</p> | <p>Knowledge Congruence = 1 iff agents have knowledge when and only when it is needful to complete their tasks. Thus, we compute the knowledge assigned to tasks via agents, and compare it with the knowledge needed for tasks.</p> <ul style="list-style-type: none"> let $KAT = (AK' * AT)$ let $d = \text{card}\{ (i,j) \mid (KAT(i,j) > 0) \neq (KT(i,j) > 0) \}$ let $d = d / (K * T)$, which normalizes d to be in $[0,1]$ <p>Then Knowledge Congruence = $1 - d$</p> |

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| Congruence, Resource | Measures the similarity between what resources are assigned to tasks via agents, and what resources are required to do tasks. Perfect congruence occurs when agents have access to resources when and only when it is needful to complete tasks. Type Graph Level Input AR:binary; AT:binary; RT:binary Output $\mathcal{R} \in [0,1]$ | Carley, 2002 | Identical to Knowledge Congruence with AR replaced by AK and KT replaced by RT. |
| Connectedness | Measures the degree to which a square network's underlying (undirected) network is connected. Type Graph Level Input N:square, symmetric Output $\mathcal{R} \in [0,1]$ | Krackhardt, 1994 | The Connectedness of a square, symmetric network is the Density of its Reachability Network. |
| Constraint, Burt | The degree to which each node in a square network is constrained from acting because of its existing links to other nodes. Type Node Level Input N:square Output $\mathcal{R} \in [0,1]$ | Burt, 1992 | This is the Constraint measure described by Equ. 2.4 on pg. 55 of Burt, 1992. Note that the matrix Z is the adjacency matrix representation of the network N. |
| Density | The ratio of the number of edges versus the maximum possible edges for a network. Type Graph Level Input N Output $\mathcal{R} \in [0,1]$ | Wasserman and Faust, 1994 (pg 101) | Let M be the adjacency matrix for the network of dimension m x n. If the network is unimodal, then m=n and M has a zero diagonal, and therefore Density = $\text{sum}(M)/(m*(m-1))$. If the network is symmetric, then Density is multiplied by two. For bimodal networks, Density = $\text{sum}(M)/(m*n)$. |
| Diameter | The maximum shortest path length between any two nodes in a unimodal network $G=(V,E)$. If there exist i,j in V such that j is not reachable from i, then V is returned. Type Graph Level Input N:square Output $Z \in [0, V]$ | Wasserman and Faust, 1994 (pg 111) | The diameter of $G=(V,E)$ is defined as: $\max \{d_G(i, j) \mid i, j \in V\}$ That is, the maximum shortest directed path between any two vertices in G. If there exists i and j such that j is not reachable from i, then V is returned. |

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| Distance Weighted Reach | <p>A generalization of graph theoretic distance, this measures the distance from a <i>set</i> of nodes in the network to all other nodes.</p> <p>Type Graph Level Input N:square, undirected Output $\mathcal{R} \in [0,1]$</p> | Borgatti, 2003 | <p>Consider a square, undirected network represented by $G=(V,E)$. let $S \subseteq V$ For any $j \notin S$, define $d_G(S, j) = \min \{d_G(i, j) \mid i \in S\}$.</p> $\sum_{j \notin S} \frac{1}{d_G(S, j)}$ <p>Then, Distance Weighted Reach = $1 - \frac{\sum_{j \notin S} \frac{1}{d_G(S, j)}}{ V - S }$,</p> |
| Diversity, Knowledge | <p>The distribution of difference in idea sharing. This is the Herfindahl-Hirshman index applied to column sums of AK.</p> <p>Type Graph Level Input AK:binary Output $\mathcal{R} \in [0,1]$</p> | | <p>This is the Herfindahl-Hirshman index (economics: sum of the squares of each firm's market share) applied to the normalized column sums of AK. This measures the degree to which knowledge is equally known.</p> <p>let $w_k = \sum_{i=1}^{ A } A(i, k)$, for $1 \leq k \leq K$</p> <p>let $W = \sum_{k=1}^{ K } w_k$</p> <p>Then Diversity = $1 - \sum_{k=1}^{ K } (w_k / W)^2$</p> |
| Diversity, Resource | <p>The distribution of difference in resource sharing. This is the Herfindahl-Hirshman index applied to column sums of AR.</p> <p>Type Graph Level Input AR:binary Output $\mathcal{R} \in [0,1]$</p> | | <p>Identical to Knowledge Diversity, with AK replaced by AR.</p> |
| Edge Count, Lateral | <p>The percentage of lateral edges in a unimodal network. Fixing a root node x, a lateral edge (i,j) is one in which the distance from x to i is the same as the distance from x to j.</p> <p>Type Graph Level Input N:square Output $\mathcal{R} \in [0,1]$</p> | Carley, 2002 | <p>Let $G=(V,E)$ be the graph representation of a unimodal network. And fix a node $x \in V$ to be the root node.</p> <p>Let $S = \{(i,j) \in E \mid d_G(x, i) = d_G(x, j)\}$</p> <p>Then Lateral Edge Count = S / E</p> |

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| Edge Count, Pooled | The percentage of pooled edges in a unimodal network. A pooled is an edge (i,j) such that there exists at least one other edge (i,k) in the network, and $k \neq j$. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Let M be the adjacency matrix representation of a unimodal network. Let $S = \{ (i,j) \mid M(i,j)=1 \wedge \text{sum}(M(:,j))>1 \}$ In other words: edge (i,j) is a pooled edge iff the in-degree of node j > 1. Then Pooled Edge Count = $ S / E $ |
| Edge Count, Reciprocal | The percentage of edges in a unimodal network that are reciprocated (also called Reciprocity). An edge (i,j) in the network is reciprocated if edge (j,i) is also in the network. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | | Let $G=(V,E)$ be the graph representation of a unimodal network. Let $S = \text{card}\{(i,j) \in E \mid i < j, (j,i) \in E \}$ Then Reciprocal Edge Count = $ S / E $ |
| Edge Count, Sequential | The percentage of edges in a unimodal network that are neither Reciprocal Edges nor Pooled Edges. Note that an edge can be both a Pooled and a Reciprocal edge. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Let $G=(V,E)$ be the graph representation of a unimodal network, and let X = set of Pooled edges of G, and let Y = set of Reciprocal edges of G. Then Sequential Edge Count = $ E-X-Y / E $ |
| Edge Count, Skip | The fraction of edges in a unimodal network that skip levels. An edge (i,j) is a skip edge if there is a path from node i to node j even after the edge (i,j) is removed. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | A skip edge in a unimodal network represented by $G=(V,E)$ is an edge $(i,j) \in E$ such that j is reachable from i in the graph $G'=(V,E \setminus (i,j))$, that is, the graph G with edge (i,j) removed. Skip Count is simply the number of such edges in G normalized to be in [0,1] by dividing by E . |
| Effective Network Size | The effective size of a node's ego network based on redundancy of ties. Type Node Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Burt, 1992 | This is the Effective Size of Network measure described by Equ. 2.2 on pg. 52 of Burt, 1992. Note that the matrix Z is the adjacency matrix representation of the network N. |

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| Efficiency | The degree to which each component in a network contains the minimum edges possible to keep it connected. Type Graph Level Input N:square, symmetric Output $\mathfrak{R} \in [0,1]$ | Krackhardt, 1994 | Let $G=(V,E)$ be the graph representation of a square, symmetric network. let n = number of components in G let c_i = number of nodes in component i let $\text{penalty} = E - V + C$ let $\text{maxPenalty} = C - V + \sum_i c_i (c_i - 1) / 2$ Then $\text{Efficiency} = 1 - \text{penalty}/\text{maxPenalty}$ |
| Exclusivity, Knowledge | Detects agents who have singular knowledge. Type Node Level Input AK:binary Output $\mathfrak{R} \in [0,1]$ | Ashworth, 2003 | The Knowledge Exclusivity Index (KEI) for agent i is defined as follows: $\sum_{j=1}^{ K } AK(i, j) * \exp(1 - \text{sum}(AK(:, j)))$ |
| Exclusivity, Resource | Detects agents who have singular resource access. Type Node Level Input AR:binary Output $\mathfrak{R} \in [0,1]$ | Ashworth, 2003 | The Resource Exclusivity Index (REI) for agent i is defined exactly as for Knowledge Based Exclusivity, but with the matrix AK replaced by AR. |
| Exclusivity, Task | Detects agents who exclusively perform tasks. Type Node Level Input AT:binary Output $\mathfrak{R} \in [0,1]$ | Ashworth, 2003 | The Task Exclusivity Index (TEI) for agent i is defined exactly as for Knowledge Based Exclusivity, but with the matrix AK replaced by AT. |
| Fragmentation | The proportion of nodes in a network that are disconnected. Type Graph Level Input N:square, undirected Output $\mathfrak{R} \in [0,1]$ | Borgatti, 2003 | Consider a square, undirected network represented by $G=(V,E)$. let $n = V $ let s_k be the number of nodes in the k^{th} component of G , $1 \leq k \leq n$ $\sum_k s_k (s_k - 1)$ Then, $\text{Fragmentation} = 1 - \frac{\sum_k s_k (s_k - 1)}{n(n - 1)}$. |
| Hierarchy | The degree to which a unimodal network exhibits a pure hierarchical structure. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Krackhardt, 1994 | Let N be a unimodal network. The Hierarchy of N is the Reciprocity of the Reachability Network for N . |

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| Interdependence | The percentage of edges in a unimodal network that are Pooled or Reciprocal. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Let $G=(V,E)$ be the graph representation of a square network. Let a = Pooled Edge Count and b = Reciprocal Edge Count of the network. Then Interdependence = $(a+b)/ E $ |
| Interlockers and Radials | Interlocker and radial nodes in a square network have a high and low Triad Count, respectively. Type Node Level Input N:square Output Binary | Carley, 2002 | Let $N=(V,E)$ be a square network. Let t_i = Triad Count for node i , $1 \leq i \leq V $. Let u = the mean of $\{t_i\}$ Let d = the variance of $\{t_i\}$ Then if $t_k \geq (u + d)$, then agent k is an <i>interlocker</i> . If $t_k \leq (u - d)$ then agent k is a <i>radial</i> . |
| Load, Knowledge | Average number of knowledge per agent. Type Graph Level Input AK:binary Output $\mathfrak{R} \in [0, K]$ | Carley, 2002 | Knowledge Load = $\text{sum}(AK) / (A)$ |
| Load, Resource | Average number of resources per agent. Type Graph Level Input AR:binary Output $\mathfrak{R} \in [0, R]$ | Carley, 2002 | Resource Load = $\text{sum}(AR) / (A)$ |
| Negotiation, Knowledge | The extent to which agents need to negotiate with each other because they lack the knowledge to complete their assigned tasks. Type Graph Level Input AT:binary; AK:binary; KT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Compute the percentage of tasks that lack at least one resource: let $\text{Need} = (AT * AK) - KT$ let $S = \{i \mid 1 \leq i \leq T , \exists j : \text{Need}(i,j) < 0\}$ Then Need for Negotiation = $ S / T $ |
| Negotiation, Resource | The extent to which agents need to negotiate with each other because they lack the resources to complete their assigned tasks. Type Graph Level Input AT:binary; AR:binary; RT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Identical to Knowledge Negotiation, replacing AK with AR, and KT with RT. |

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| Network Centralization, Betweenness | Network centralization based on the betweenness score for each node in a square network. This measure is defined for directed and undirected networks. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Freeman, 1979 | Let $G=(V,E)$ represent the square network, and let $n = V $ let $d_i =$ Betweenness Centrality of node i let $\bar{d} = \max\{d_i 1 \leq i \leq n\}$ Then Network Betweenness Cent. = $\left(\sum_{1 \leq i \leq n} \bar{d} - d_i \right) / (n-1)$. |
| Network Centralization, Closeness | Network centralization based on the closeness centrality of each node in a square network. This is defined only for connected, undirected networks. Type Graph Level Input N:square, symmetric, connected Output $\mathfrak{R} \in [0,1]$ | Freeman, 1979 | Let $G=(V,E)$ represent the square network, and let $n = V $ let $d_i =$ Closeness Centrality of node i let $\bar{d} = \max\{d_i 1 \leq i \leq n\}$ Then Network Closeness Cent. = $\left(\sum_{1 \leq i \leq n} \bar{d} - d_i \right) / ((n-2)(n-1)/(2n-3))$. |
| Network Centralization, Column Degree | A centralization based on the degree of the column nodes of a network. Type Graph Level Input N Output $\mathfrak{R} \in [0,1]$ | NetStat | Let N be a network with n column nodes. let $d_j =$ degree of column node j , $1 \leq j \leq n$ let $\bar{d} = \max\{d_j 1 \leq j \leq n\}$ Then Column Degree Network Centralization = $\left(\sum_{1 \leq j \leq n} \bar{d} - d_j \right) / (n)$. |
| Network Centralization, In Degree | A centralization of a square network based on the In-Degree Centrality of each node. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | NetStat | Let N be a unimodal network with n nodes. let $d_i =$ In Degree Centrality of node i let $\bar{d} = \max\{d_i 1 \leq i \leq n\}$ Then In Degree Network Centralization = $\left(\sum_{1 \leq i \leq n} \bar{d} - d_i \right) / D$, where $D = (n-2)$ if N is undirected, and $(n-1)$ otherwise. |

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| Network Centralization, Out Degree | A centralization of a square network based on the Out-Degree Centrality of each node. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | NetStat | Let N be a unimodal network with n nodes. let $d_i =$ Out Degree Centrality of node i let $\bar{d} = \max\{d_i 1 \leq i \leq n\}$ Then Out Degree Network Centralization = $\left(\sum_{1 \leq i \leq n} \bar{d} - d_i \right) / D$, where D = (n-2) if N is undirected, and (n-1) otherwise. |
| Network Centralization, Row Degree | A centralization based on the degree of the row nodes in a network. Type Graph Level Input N Output $\mathfrak{R} \in [0,1]$ | NetStat | Let N be a network with n row nodes. let $d_j =$ degree of row node j, $1 \leq j \leq n$ let $\bar{d} = \max\{d_j 1 \leq j \leq n\}$ Then Row Degree Network Centralization = $\left(\sum_{1 \leq j \leq n} \bar{d} - d_j \right) / (n)$. |
| Network Centralization, Total Degree | A centralization of a square network based on total degree centrality of each node. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Freeman, 1979 | Let N be a unimodal network with n nodes. let $d_i =$ Total Degree Centrality of node i let $\bar{d} = \max\{d_i 1 \leq i \leq n\}$ Then Total Degree Network Centralization = $\left(\sum_{1 \leq i \leq n} \bar{d} - d_i \right) / (n - 2)$. |
| Network Levels | The Network Level of a square network is the maximum Node Level of its nodes. Type Graph Level Input N:square Output $Z \in [0, V - 1]$ | NetStat | Let $G=(V,E)$ be the graph representation of a square network. Then the Levels of $G = \max \{ d_G(i, j) \mid i, j \in V; j \text{ reachable from } i \text{ in } G \}$ |
| Node Level | The Node Level for a node v in a square network is the longest shortest path from v to every node v can reach. If v cannot reach any node, then its level is 0. Type Node Level Input N:square Output $Z \in [0, V - 1]$ | Carley, 2002 | Let $G=(V,E)$ be the graph representation of a square network and fix a node v. Node Level for v = $\max \{ d_G(v, j) \mid j \in V; j \text{ reachable from } v \text{ in } G \}$; if v cannot reach any nodes, then its level is 0. |

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| Omega, Knowledge | <p>The degree to which agents reuse knowledge while doing their tasks.</p> <p>Type Graph Level</p> <p>Input AT:binary; KT:binary; T:binary</p> <p>Output $\mathfrak{R} \in [0,1]$</p> | Carley, Dekker, and Krackhardt 2000 | <p>Let $TAT = TA*TA'$</p> <p>Let $N = ((T' @TAT)*KT')@KT'$</p> <p>Then Knowledge Based Omega = $\text{sum}(N)/\text{sum}(KT)$</p> |
| Omega, Resource | <p>The degree to which agents reuse resources while doing their tasks.</p> <p>Type Graph Level</p> <p>Input AT:binary; RT:binary; T:binary</p> <p>Output $\mathfrak{R} \in [0,1]$</p> | Carley, Dekker, and Krackhardt 2000 | Identical to Knowledge Based Omega, replacing KT with RT. |
| Performance as Accuracy | <p>Measures how accurately agents can perform their assigned tasks based on their access to knowledge and resources.</p> <p>Type Graph Level</p> <p>Input AT:binary; AK:binary; AR:binary; KT:binary; RT:binary</p> <p>Output $\mathfrak{R} \in [0,1]$</p> | Carley, 2002 | <p>Accuracy is computed based on the binary classification problem. It is computed in one of two ways:</p> <p>(1) Knowledge based: Let b be a binary string of length K, let $N=KT'$, and let $S=AK$. Fix a task t.</p> <p>let $\text{answer} = (\sum_{1 \leq k \leq K } N(t,k)b_k / \sum_{1 \leq k \leq K } N(t,k) > .5)$, which is the correct classification of b with respect to task t. Now, let $I=\{ i \mid AT(i,t)=1 \}$.</p> <p>let $\text{answer}(i) = (\sum_{1 \leq k \leq K } N(t,k)S(i,k)b_k / \sum_{1 \leq k \leq K } N(t,k)S(i,k) > .5)$, $i \in I$.</p> <p>This is agent i's classification of b with respect to t.</p> <p>The group of agents classify b using majority voting. That is, let</p> $\text{group_answer} = (\frac{1}{ I } \sum_{i \in I} \text{answer}(i) > .5).$ <p>Then, if $\text{group_answer} = \text{answer}$, then the group was accurate, otherwise not.</p> <p>This is repeated multiple times for each task, and across all tasks. The percentage correct is Performance as Accuracy.</p> <p>(2) Resource based: let $N=RT'$ and $S=AR$ in the analysis of case (1).</p> <p>If the network has the knowledge and resource graphs to perform both cases, then Performance as Accuracy is the average of the two.</p> |

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| Personnel Cost | Total number of people reporting to an agent, plus its total knowledge, resources, and tasks. Type Node Level Input A:binary; AK:binary; AR:binary; AT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2003 | Personnel Cost for agent i is defined as follows: Let $d = \text{sum}(A(:,i)) + \text{sum}(AK(i,:)) + \text{sum}(AR(i,:)) + \text{sum}(AT(i,:))$ The value is then normalized to be in [0,1]: Let $d = d / ((A -1) + K + R + T)$ The Personnel Cost for agent i is d. |
| Potential Workload, Knowledge | Maximum knowledge an agent could use to do tasks if it were assigned to all tasks. Type Node Level Input AK:binary; KT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Potential Knowledge Workload for agent i = $\text{sum}((AK*KT)(i,:))/\text{sum}(KT)$ |
| Potential Workload, Resource | Maximum resources an agent could use to do tasks if it were assigned to all tasks. Type Node Level Input AR:binary; RT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Potential Resource Workload for agent i is identical to Potential Knowledge Workload, replacing AK with AR, and KT with RT. |
| Redundancy, Access | Average number of redundant agents per resource. An agent is redundant if there is already an agent that has access to the resource. Type Graph Level Input AR:binary Output $\mathfrak{R} \in [0, (A - 1) * R]$ | Carley, 2002 | This is the Column Redundancy of matrix AR. |
| Redundancy, Assignment | Average number of redundant agents assigned to tasks. An agent is redundant if there is already an agent assigned to the task. Type Graph Level Input AT Output $\mathfrak{R} \in [0, (A - 1) * T]$ | Carley, 2002 | This is the Column Redundancy of matrix AT. |

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| Redundancy, Column | The mean number of column node edges in excess of one. Type Graph Level Input N of dimension m x n Output $\mathfrak{R} \in [0, (m-1) * n]$ | Netstat | Let M be the matrix representation for a network N of dimension m x n. let $d_j = \max\{0, \text{sum}(M(:, j)) - 1\}$, for $1 \leq j \leq n$; this is the number of column entries in excess of one for column j. Then Column Redundancy = $\left(\sum_{j=1}^n d_j \right) / n$ |
| Redundancy, Knowledge | Average number of redundant agents per knowledge. An agent is redundant if there is already an agent that has the knowledge. Type Graph Level Input AK Output $\mathfrak{R} \in [0, (A - 1) * K]$ | Carley, 2002 | This is the Column Redundancy of matrix AK. |
| Redundancy, Resource | Average number of redundant resources assigned to tasks. A resource is redundant if there is already a resource assigned to the task. Type Graph Level Input RT:binary Output $\mathfrak{R} \in [0, (R - 1) * T]$ | Carley, 2002 | This is the Column Redundancy of matrix RT. |
| Redundancy, Row | The mean number of row node edges in excess of one. Type Graph Level Input N of dimension m x n Output $\mathfrak{R} \in [0, (n-1) * m]$ | Netstat | Let M be the matrix representation for a network N of dimension m x n. let $d_i = \max\{0, \text{sum}(M(j, :)) - 1\}$, for $1 \leq i \leq m$; this is the number of column entries in excess of one for row i. Then Row Redundancy = $\left(\sum_{j=1}^m d_j \right) / m$ |
| Relative Expertise | The degree of dissimilarity between agents based on shared knowledge. Each agent computes to what degree the other agents know what they do not know. Type Node Level Input AK:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | The Relative Expertise matrix (RE) is defined as follows: RE(i,i) = 0 RE(i,j) = ($\sim AK * AK'$) = # knowledge that j knows that i does not know Finally, normalize RE by its row sums: RE(i,:) /= sum(RE(i,:)) The Relative Expertise for agent i = $\left(\sum_{\substack{j=1 \\ j \neq i}}^{ A } RE(i, j) \right) / (A - 1)$, that is, the average of the non-diagonal elements of row i of RE. |

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| Relative Similarity | The degree of similarity between two agents based on shared knowledge. Each agent computes to what degree the other agents know what they know. Type Node Level Input AK: binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Let $M = AK * AK'$ Let $w(i) = \text{sum}(M(i,:), 1 \leq i \leq A)$ Then Relative Similarity (RS) between agents i and j is $RS(i,j) = M(i,j)/w(i)$. The Relative Similarity for an agent i = $\left(\sum_{\substack{j=1 \\ j \neq i}}^{ A } RS(i, j) \right) / (A - 1)$, that is, the average of the non-diagonal elements of row i of RS. |
| Span of Control | The average number of out edges per node with non-zero out degrees. Type Graph Level Input N:square Output $\mathfrak{R} \in [0, V - 1]$ | Carley, 2002 | let S = set of nodes in V that have positive out-degree let $K = \sum_{i \in S} outDegree(i)$ Then Span of Control = $K / S $ |
| Speed, Average | The average shortest path length between node pairs (i,j) where there is a path in the network from i to j. If there are no such pairs, then Average Speed is zero. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | let $G=(V,E)$ be the graph representation of a square network. let $D=\{(i,j) \mid i,j \in V, j \text{ reachable from } i \text{ in } G\}$ Then Average Speed = $\left(\sum_{(i,j) \in D} d_G(i, j) \right) / D $ |
| Speed, Minimum | The maximum shortest path length between node pairs (i,j) where there is a path in the network from i to j. If there are no such pairs, then Minimum Speed is zero. Type Graph Level Input A Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Minimum Speed = $1 / (\text{Levels for the Communication Network})$ |
| Task Completion, Knowledge Based | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite knowledge to do the tasks. Type Graph Level Input AK:binary; AT:binary; KT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Find the tasks that cannot be completed because the agents assigned to the tasks lack necessary knowledge: let $Need = [(AT * AK) - KT]$ let $S = \{ i \mid 1 \leq i \leq T , \exists j : Need(i,j) < 0 \}$ Knowledge Based Task Completion is the percentage of tasks that could be completed = $(T - S) / T $ |

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| Task Completion, Overall | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite knowledge and resources to do the tasks. Type Graph Level Input AR:binary; AT:binary; RT:binary; AK:binary, KT:binary Output $\mathcal{R} \in [0,1]$ | Carley, 2002 | This is the average of Knowledge Based Task Completion and Resource Based Task Completion. If one of the two could not be computed, then the other is returned. |
| Task Completion, Resource Based | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite resources to do the tasks. Type Graph Level Input AR:binary; AT:binary; RT:binary Output $\mathcal{R} \in [0,1]$ | Carley, 2002 | Find the tasks that cannot be completed because the agents assigned to the tasks lack necessary resources. Defined identically as Knowledge Based Task Completion, replacing matrix AK with AR and matrix KT with RT. |
| Transitivity | The percentage of edge pairs $\{(i,j), (j,k)\}$ in the network such that (i,k) is also an edge in the network. Type Graph Level Input N:square Output $\mathcal{R} \in [0,1]$ | NetStat | Let $G = (V,E)$ be the graph representation of the square network. let $I = \{(i,j,k) \in V^3 \mid i,j,k \text{ distinct}\}$ let $Potential = \{(i,j,k) \in I \mid (i,j) \in E, \text{ and } (j,k) \in E\}$ let $Complete = \{(i,j,k) \in Potential \mid (i,k) \in E\}$ Then $Transitivity = Complete / Potential $ |
| Triad Count | The number of triads centered at each node in a square network. Type Node Level Input N:square of dimension $ V $ Output $Z \in [0, (V - 1)(V - 2)]$ | NetStat | Let $G=(V,E)$ represent a square network. And let Triad be a matrix of dimension $ V \times V $. $Triad(i,i) = 0$ $Triad(i,j) = \text{card}\{k \mid k \neq i, k \neq j; A(i,j) \wedge A(i,k) \wedge A(k,j)\}, i \neq j$ Then the Triad count for agent $i = \text{sum}(Triad(i,:))$ |
| Under Supply, Knowledge | The extent to which the knowledge needed to do tasks are unavailable in the entire organization. Type Graph Level Input AK:binary; AT:binary; KT:binary Output $\mathcal{R} \in [0,1]$ | Carley, 2002 | Compute the average number of needed knowledge per task: let $Need = (AT^*AK) - KT$ let $TaskNeed(i) = \text{card}\{j \mid Need(i,j) < 0\}, \text{ for } 1 \leq i \leq T $ Then UnderSupply is $\text{sum}(TaskNeed) / T $ |

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| Under Supply, Resource | The extent to which the resources needed to do tasks are unavailable in the entire organization. Type Graph Level Input AR:binary; AT:binary; RT:binary Output $\mathfrak{R} \in [0,1]$ | Carley, 2002 | Under Resource Supply is identical to Under Knowledge Supply, replacing AK with AR, and KT with RT. |
| Upper Boundedness | The degree to which pairs of agents have a common ancestor. Type Graph Level Input N:square Output $\mathfrak{R} \in [0,1]$ | Krackhardt, 1994 | |
| Weak Boundary Spanner | A node which if removed from a network creates a new component. Type Node Level Input N:square, symmetric Output Binary | Cormen, Leiserson, Rivest, Stein, 2001 p.558 | A Weak Boundary Spanner is an <i>articulation point</i> of N, as defined in the referenced book. |

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Appendix B – ORA Risk Category Measures

| Metric | Meaning | Level | Risk | Data Needs |
|-------------------------------|--|--------------|---------------------|-------------------|
| Access Index, Knowledge Based | Boolean value which is true if an agent is the only agent who knows a piece of knowledge and who is known by exactly one other agent. The one agent known also has its KAI set to one. | Node | Critical Employee | Multi-cell |
| Access Index, Resource Based | Boolean value which is true if an agent is the only agent with access to a resource and who is known by exactly one other agent. The one agent known also has its RAI set to one. | Node | Critical Employee | Multi-cell |
| Actual Workload, Knowledge | The knowledge an agent uses to perform the tasks to which it is assigned. | Node | Resource Allocation | Multi-cell |
| Actual Workload, Resource | The resources an agent uses to perform the tasks to which it is assigned. | Node | Resource Allocation | Multi-cell |
| Average Distance | The average shortest path length between nodes, excluding infinite distances. | Graph | Communication | Single-cell |
| Centrality, Betweenness | The Betweenness Centrality of node v in a network is defined as: across all node pairs that have a shortest path containing v , the percentage that pass through v . This is defined for directed networks. | Node | Communication | Single-cell |
| Centrality, Closeness | The average closeness of a node to the other nodes in a network. Loosely, Closeness is the inverse of the average distance in the network between the node and all other nodes. This is defined for directed networks. | Node | Communication | Single-cell |
| Centrality, Eigenvector | Calculates the eigenvector of the largest positive eigenvalue of the adjacency matrix representation of a square network. | Node | Communication | Single-cell |
| Centrality, In Degree | The In Degree Centrality of a node in a unimodal network is its normalized in-degree. | Node | Communication | Single-cell |
| Centrality, Information | Calculate the Stephenson and Zelen information centrality measure for each node. | Node | Communication | Single-cell |

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| Centrality, Inverse Closeness | The average closeness of a node to the other nodes in a network. Inverse Closeness is the sum of the inverse distances between a node and all other nodes. This is defined for directed networks. | Node | Communication | Single-cell |
| Centrality, Out Degree | The Out Degree Centrality of a node in a square network is its normalized out-degree. | Node | Communication | Single-cell |
| Centrality, Total Degree | The Total Degree Centrality of a node in a square network is its normalized in plus out degree. | Node | Communication | Single-cell |
| Clustering Coefficient, Watts-Strogatz | Measures the degree of clustering in a network by averaging the clustering coefficient of each node i , defined as the ratio of the number of triangles connected to i to the number of triples centered at i . | Graph | Communication | Single-cell |
| Cognitive Load | Measures the total amount of effort expended by each agent to do its tasks. | Node | Critical Employee | Multi-cell |
| Communication | Measures the communication need of agents to complete their assigned tasks. | Node | Communication | Multi-cell |
| Component Count, Strong | The number of strongly connected components in a network. | Graph | Communication | Single-cell |
| Component Count, Weak | The number of weakly connected components in a network. | Graph | Communication | Single-cell |
| Congruence, Communication | Measures to what extent the agents communicate when and only when it is needful to complete tasks. Perfect congruence requires reciprocal communication. | Graph | Communication | Multi-cell |
| Congruence, Knowledge | Measures the similarity between what knowledge is assigned to tasks via agents, and what knowledge is required to do tasks. Perfect congruence occurs when agents have knowledge when and only when it is needful to complete tasks. | Graph | Resource Allocation, Task | Multi-cell |
| Congruence, Resource | Measures the similarity between what resources are assigned to tasks via agents, and what resources are required to do tasks. Perfect congruence occurs when agents have access to resources when and only when it is needful to complete tasks. | Graph | Resource Allocation, Task | Multi-cell |

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| Connectedness | Measures the degree to which a square network's underlying (undirected) network is connected. | Graph | Communication | Single-cell |
| Constraint, Burt | The degree to which each node in a square network is constrained from acting because of its existing links to other nodes. | Node | Critical Employee, Redundancy, Communication | Single-cell |
| Density | The ratio of the number of edges versus the maximum possible edges for a network. | Graph | | Single-cell |
| Diameter | The maximum shortest path length between any two nodes in a unimodal network $G=(V,E)$. If there exist i,j in V such that j is not reachable from i , then $ V $ is returned. | Graph | Communication | Single-cell |
| Distance Weighted Reach | A generalization of graph theoretic distance, this measures the distance from a <i>set</i> of nodes in the network to all other nodes. | Graph | Communication, Critical Employee | Single-cell |
| Diversity, Knowledge | The distribution of difference in idea sharing. This is the Herfindahl-Hirshman index applied to column sums of AK. | Graph | Resource Allocation | Single-cell |
| Diversity, Resource | The distribution of difference in resource sharing. This is the Herfindahl-Hirshman index applied to column sums of AR. | Graph | Resource Allocation | Single-cell |
| Edge Count, Lateral | The percentage of lateral edges in a unimodal network. Fixing a root node x , a lateral edge (i,j) is one in which the distance from x to i is the same as the distance from x to j . | Graph | Communication | Single-cell |
| Edge Count, Pooled | The percentage of pooled edges in a unimodal network. A pooled is an edge (i,j) such that there exists at least one other edge (i,k) in the network, and $k \neq j$. | Graph | Task | Single-cell |
| Edge Count, Reciprocal | The percentage of edges in a unimodal network that are reciprocated (also called Reciprocity). An edge (i,j) in the network is reciprocated if edge (j,i) is also in the network. | Graph | Task | Single-cell |
| Edge Count, Sequential | The percentage of edges in a unimodal network that are neither Reciprocal Edges nor Pooled Edges. Note that an edge can be both a Pooled and a Reciprocal edge. | Graph | Task | Single-cell |

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| Edge Count, Skip | The fraction of edges in a unimodal network that skip levels. An edge (i,j) is a skip edge if there is a path from node i to node j even after the edge (i,j) is removed. | Graph | Communication | Single-cell |
| Effective Network Size | The effective size of a node's ego network based on redundancy of ties. | Node | Critical Employee, Redundancy, Communication | Single-cell |
| Exclusivity, Knowledge | Detects agents who have singular knowledge. | Node | Critical Employee | Single-cell |
| Exclusivity, Resource | Detects agents who have singular resource access. | Node | Critical Employee | Single-cell |
| Exclusivity, Task | Detects agents who exclusively perform tasks. | Node | Critical Employee Performance | Single-cell |
| Fragmentation | The proportion of nodes in a network that are disconnected | Graph | Critical Employee, Communication, Personnel Interaction | Single-cell |
| Hierarchy | The degree to which a unimodal network exhibits a pure hierarchical structure. | Graph | Communication | Single-cell |
| Interdependence | The percentage of edges in a unimodal network that are Pooled or Reciprocal. | Graph | Task | Single-cell |
| Interlockers and Radials | Interlocker and radial nodes in a square network have a high and low Triad Count, respectively. | Graph | Critical Employee | Single-cell |
| Load, Knowledge | Average number of knowledge per agent. | Graph | Resource Allocation, Redundancy | Single-cell |
| Load, Resource | Average number of resources per agent. | Graph | Resource Allocation, Redundancy | Single-cell |
| Negotiation, Knowledge | The extent to which agents need to negotiate with each other because they lack the knowledge to complete their assigned tasks. | Graph | Resource Allocation, Task | Multi-cell |
| Negotiation, Resource | The extent to which agents need to negotiate with each other because they lack the resources to complete their assigned tasks. | Graph | Resource Allocation, Task | Multi-cell |
| Network Centralization, Betweenness | Network centralization based on the betweenness score for each node in a square network. This measure is defined for directed and undirected networks. | Graph | Communication | Single-cell |
| Network Centralization, Closeness | Network centralization based on the closeness centrality of each node in a square network. This is defined only for connected, undirected networks. | Graph | Communication | Single-cell |

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| Network Centralization, Column Degree | A centralization based on the degree of the column nodes of a network. | Graph | | Single-cell |
| Network Centralization, In Degree | A centralization of a square network based on the In-Degree Centrality of each node. | Graph | Communication | Single-cell |
| Network Centralization, Out Degree | A centralization of a square network based on the Out-Degree Centrality of each node. | Graph | Communication | Single-cell |
| Network Centralization, Row Degree | A centralization based on the degree of the row nodes in a network. | Graph | | Single-cell |
| Network Centralization, Total Degree | A centralization of a square network based on total degree centrality of each node. | Graph | Communication | Single-cell |
| Network Levels | The Network Level of a square network is the maximum Node Level of its nodes. | Graph | Communication | Single-cell |
| Node Level | The Node Level for a node v in a square network is the longest shortest path from v to every node v can reach. If v cannot reach any node, then its level is 0. | Graph | Communication | Single-cell |
| Omega, Knowledge | The degree to which agents reuse knowledge while doing their tasks. | Graph | Task, Performance | Multi-cell |
| Omega, Resource | The degree to which agents reuse resources while doing their tasks. | Graph | Task, Performance | Multi-cell |
| Performance as Accuracy | Measures how accurately agents can perform their assigned tasks based on their access to knowledge and resources. | Graph | Task, Performance | Multi-cell |
| Personnel Cost | Total number of people reporting to an agent, plus its total knowledge, resources, and tasks. | Node | | Multi-cell |
| Potential Workload, Knowledge | Maximum knowledge an agent could use to do tasks if it were assigned to all tasks. | Node | Resource Allocation | Multi-cell |
| Potential Workload, Resource | Maximum resources an agent could use to do tasks if it were assigned to all tasks. | Node | Resource Allocation | Multi-cell |
| Redundancy, Access | Average number of redundant agents per resource. An agent is redundant if there is already an agent that has access to the resource. | Graph | Resource Allocation, Redundancy | Multi-cell |
| Redundancy, Assignment | Average number of redundant agents assigned to tasks. An agent is redundant if there is already an agent assigned to the task. | Graph | Resource Allocation, Redundancy | Multi-cell |
| Redundancy, Column | The mean number of column node edges in excess of one. | Graph | | Single-cell |

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| Redundancy, Knowledge | Average number of redundant agents per knowledge. An agent is redundant if there is already an agent that has the knowledge. | Graph | Resource Allocation, Redundancy, Task | Single-cell |
| Redundancy, Resource | Average number of redundant resources assigned to tasks. A resource is redundant if there is already a resource assigned to the task. | Graph | Resource Allocation, Redundancy, Task | Single-cell |
| Redundancy, Row | The mean number of row node edges in excess of one. | Graph | | Single-cell |
| Relative Expertise | The degree of dissimilarity between agents based on shared knowledge. Each agent computes to what degree the other agents know what they do not know. | Node | Personnel Interaction | Single-cell |
| Relative Similarity | The degree of similarity between two agents based on shared knowledge. Each agent computes to what degree the other agents know what they know. | Node | Personnel Interaction | Single-cell |
| Span of Control | The average number of out edges per node with non-zero out degrees. | Graph | Communication | Single-cell |
| Speed, Average | The average shortest path length between node pairs (i,j) where there is a path in the network from i to j. If there are no such pairs, then Average Speed is zero. | Graph | Communication | Single-cell |
| Speed, Minimum | The maximum shortest path length between node pairs (i,j) where there is a path in the network from i to j. If there are no such pairs, then Minimum Speed is zero. | Graph | Communication | Single-cell |
| Task Completion, Knowledge Based | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite knowledge to do the tasks. | Graph | Performance | Multi-cell |
| Task Completion, Overall | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite knowledge and resources to do the tasks. | Graph | Performance | Multi-cell |
| Task Completion, Resource Based | The percentage of tasks that can be completed by the agents assigned to them, based solely on whether the agents have the requisite resources to do the tasks. | Graph | Performance | Multi-cell |

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| Transitivity | The percentage of edge pairs $\{(i,j), (j,k)\}$ in the network such that (i,k) is also an edge in the network. | Graph | Communication, Task | Single-cell |
| Triad Count | The number of triads centered at each node in a square network. | Node | Communication | Single-cell |
| Under Supply, Knowledge | The extent to which the knowledge needed to do tasks are unavailable in the entire organization. | Graph | Resource Allocation, Task | Multi-cell |
| Under Supply, Resource | The extent to which the resources needed to do tasks are unavailable in the entire organization. | Graph | Resource Allocation, Task | Multi-cell |
| Upper Boundedness | The degree to which pairs of agents have a common ancestor. | Graph | Communication | Single-cell |
| Weak Boundary Spanner | A node which if removed from a network creates a new component. | Node | Critical Employee | Single-cell |